

Math 512 Syllabus  
Spring 2019, LIU Post

Week	Class Date	Material
1	1/28	ISBN, error-detecting codes HW: Exercises 1.1, 1.3, 1.5 Find the multiplicative inverse for all non-zero elements of $\mathbb{F}_{11}$ Show that ISBN-13 need not detect adjacent swaps
2	2/4	Error probability, Repetition Codes, Hamming square code HW: Exercises 1.7-1.9, 1.14-1.16 (linear part optional) Calculate the probabilities of transmitting strings error-free using the $[3, 1]$ -repetition and the Hamming Square codes. (You may assume both correct 1 error.) For specific values of $p$ , compare with sending strings with no encoding.
3	2/11	Linear Algebra over finite fields, the beginning of Linear Codes (§2.1, §2.3-2.6). HW: Problem written on board (list all codewords), and #1 and #10 from “Homework 3” handout.
4	2/19	Hamming $[7, 4]$ code (§1.7) Linear Codes (§2.1, §2.3-2.6) HW: 4, 5c, 6, and 7 from “Homework 3” handout.
5	2/25	Generator, Parity matrices Homework 5 handout
6	3/4	<b>Quiz 1</b> Weights, distances, and detection/correction. HW: Homework 6 handout
	3/11	Spring Break
7	3/18	Review, Probabilities HW: Updated HW 6 handout
8	3/25	<b>Quiz 2</b> , Polynomial Intro HW: Homework 7 handout
9	4/1	Galois Fields HW: Homework 8 handout
10	4/8	Polynomial Codes HW: Homework 9 handout
11	4/15	Cyclic Codes HW: Homework 10 handout
12	4/22	<b>Quiz 3</b>
13	4/29	Reed–Solomon Codes
	5/6	<b>Presentations of Final Projects</b>

# Linear algebra info - MTH 512

In a linear algebra course, one is primarily concerned with linear maps between vector spaces. A map is linear if it is compatible with both vector addition and scalar multiplication.

**Definition 0.1.** A vector space over a field  $\mathbb{F}$  is a set  $V$  equipped with binary operations  $+$  (vector addition) and  $\cdot$  (scalar multiplication)

$$V \times V \xrightarrow{+} V, \quad \mathbb{F} \times V \xrightarrow{\cdot} V$$

satisfying the following axioms (for all  $\mathbf{v}, \mathbf{w}, \mathbf{x} \in V$  and  $r, s \in \mathbb{F}$ ):

1.  $\mathbf{v} + \mathbf{w} = \mathbf{w} + \mathbf{v}$  (addition commutative)
2.  $(\mathbf{v} + \mathbf{w}) + \mathbf{x} = \mathbf{v} + (\mathbf{w} + \mathbf{x})$  (addition associative)
3.  $\exists \mathbf{0} \in V$  such that  $\mathbf{v} + \mathbf{0} = \mathbf{v}$  for all  $\mathbf{v}$  (additive identity)
4.  $\forall \mathbf{v} \in V, \exists (-\mathbf{v}) \in V$  such that  $\mathbf{v} + (-\mathbf{v}) = \mathbf{0}$  (additive inverse)
5.  $r(\mathbf{v} + \mathbf{w}) = r\mathbf{v} + r\mathbf{w}$  (distributive)
6.  $(r + s)\mathbf{v} = r\mathbf{v} + s\mathbf{v}$  (distributive)
7.  $r(s\mathbf{v}) = (rs)\mathbf{v}$  (scalar associative)
8.  $1\mathbf{v} = \mathbf{v}$  (scalar identity)

**Example 0.2.**  $\mathbb{F}_p^n = \{(x_1, \dots, x_n) \mid x_i \in \mathbb{F}_p\}$ .

More specifically, consider the vectors  $(0, 1, 2, 0), (1, 0, 2, 1) \in \mathbb{F}_3^4$ . Then

$$(0, 1, 2, 0) + (1, 0, 2, 1) = (1, 0, 1, 1), \quad 2(0, 1, 2, 0) = (0, 2, 1, 0), \quad 0(0, 1, 2, 0) = (0, 0, 0, 0).$$

**Definition 0.3.** Let  $V$  be a vector space and  $W \subset V$  a subset. Then  $W$  is a *subspace* (or vector subspace or linear subspace) if

$$\mathbf{u}, \mathbf{v} \in W \quad \Rightarrow \quad r\mathbf{u} + s\mathbf{v} \in W.$$

**Definition 0.4.** Let  $V, W$  be vector spaces. A map  $T : V \rightarrow W$  is *linear* if

$$T(r\mathbf{v} + s\mathbf{w}) = rT(\mathbf{v}) + sT(\mathbf{w})$$

for all  $\mathbf{v}, \mathbf{w} \in V$  and  $r, s \in \mathbb{F}$ . A linear map  $T$  is *injective* if it is one-to-one, *surjective* if it is onto, and an *isomorphism* if it is a bijection.

**Definition 0.5.** Given a linear map  $T : V \rightarrow W$ ,

$$\text{Kernel} \quad \text{Ker}(T) := \{\mathbf{v} \in V \mid T(\mathbf{v}) = \mathbf{0} \in W\} \subseteq V,$$

$$\text{Image} \quad \text{Image}(T) := \{T(\mathbf{v}) \in W \mid \mathbf{v} \in V\} \subseteq W.$$

**Proposition 0.6.** *The Kernel and Image of a linear map are vector subspaces.*

**Definition 0.7.** Let  $V$  be a vector space over  $F$ . A finite collection of vectors  $B = \{\mathbf{v}_1, \dots, \mathbf{v}_n\} \subseteq V$  is a **basis** of  $V$  if the induced linear map

$$F^n \longrightarrow V$$

$$(r_1, \dots, r_n) \longmapsto r_1\mathbf{v}_1 + r_2\mathbf{v}_2 + \dots + r_n\mathbf{v}_n$$

is an *isomorphism*. In such a case, we say the dimension of  $V$  is  $\dim(V) = n$ .

**Theorem 0.8** (Rank-Nullity). *If  $T : V \rightarrow W$  is a linear map (and  $V$  is finite-dimensional), then*

$$\dim \text{Ker}(T) + \dim \text{Image}(T) = \dim V.$$

Math 512 - "Homework 3"  
Due February 19, 2019

1. Use row reduction/Gaussian elimination to solve the following system of linear equations over  $\mathbb{R}$ . Also, solve them over  $\mathbb{F}_5$ .

$$\begin{cases} x + 3y & = 0 \\ 3x + y + 2z & = 0 \end{cases}$$

2. Let

$$C = \left\{ (x_1, \dots, x_{10}) \in \mathbb{F}_{11}^{10} \mid \sum_{i=1}^{10} ix_i = 0 \right\} \subset \mathbb{F}_{11}^{10},$$

and let  $ISBN \subset \mathbb{F}_{11}^{10}$  be the subset of numbers which are ISBN numbers for some published book.

- (a) What is the relationship between  $C$  and  $ISBN$ ?  
(b) Determine whether  $C$  and/or  $ISBN$  are linear codes.
3. (a) Create your own example of a linear code. Explain/show why it is linear.  
(b) Create your own example of a non-linear code. Explain/show why it is non-linear.
4. Consider the  $[6, 3]$  linear code given on p. 12 of "Error-correcting codes" and discussed in class.  
(a) Write the generator matrix  $G$  for this encoding.  
(b) Use the generator matrix to send the message 101.  
(c) You receive the message 011010. Correct this if necessary/possible.  
(d) Construct the parity matrix  $H$ . Try to do this by writing the equations that any valid codeword must satisfy.  
(e) Using  $H$ , determine whether the following are valid codewords: 101101, 011010, 100011.  
(f) Calculate the matrix product  $GH^T$ .
5. This problem concerns Hamming's  $[7, 4]$ -code.  
(a) Write the parity matrix  $H$ , and use this to write the generator matrix  $G$ .  
(b) Write the digits of 3.14 as 4-bit binary numbers, and encode each of them.  
(c) You receive the following message: 1001011, 0101111, 1101001, 1110010. Correct and decode the message.  
(d) List all codewords of Hamming's  $[7, 4]$ -code.
6. Consider the  $q$ -ary repetition code code of type  $[6, 2]$ , given by taking 2 elements of  $\mathbb{F}_q$  and repeating each of them 3 times.  
(a) Write the generator matrix  $G$  and a parity matrix  $H$ .  
(b) Calculate  $GH^T$ .

7. Let  $C$  be a binary linear code with generator matrix

$$G = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 \end{bmatrix}.$$

List all the codewords in  $C$ . (You can do this by encoding all vectors in  $\mathbb{F}_2^2$ .)

8. (optional) Consider working over the field  $\mathbb{F}_5$ . Which of the following two matrices would be a valid generating matrix  $G$ ? Explain your answer. What problem would one of them cause?

$$G_1 = \begin{bmatrix} 1 & 3 & 0 \\ 3 & 1 & 2 \end{bmatrix}, \quad G_2 = \begin{bmatrix} 1 & 3 \\ 3 & 1 \\ 0 & 2 \end{bmatrix}$$

9. Let  $G_{kn}$  be an  $k \times n$  matrix with entries in  $\mathbb{F}_q$ . Show that the set of vectors  $y \in \mathbb{F}_q^n$ , satisfying  $y = xG$  for some  $x \in \mathbb{F}_q^k$ , determines a linear code.  
(Hint: We need to show that  $C = \{xG \in \mathbb{F}_q^n \mid x \in \mathbb{F}_q^k\}$  is a linear subspace of  $\mathbb{F}_q^n$ . To do this: assume  $y_1, y_2 \in C$ , and then prove  $ay_1 + by_2 \in C$ .)
10. Let  $H$  be matrix with entries in  $\mathbb{F}_q$ . Show that the set of vectors  $y$  satisfying  $yH^T = 0$  determines a linear code. In order for this to make sense, what is the relationship between the length of  $y$  and the dimensions of  $H$ ?
11. For your edification, read the brief story of transmission of photographs from deep-space, taken from Hill's "A first course in coding theory," [which you can see by clicking here](#). (You don't have to do the Exercises.)

Math 512 - Homework 5  
Due March 4, 2019

Next week's quiz will be, essentially, #1 parts a,b,c.

1. This problem concerns Hamming's  $[7, 4]$ -code.
  - (a) Write the parity matrix  $H$ , and use this to write the generator matrix  $G$ .
  - (b) Write the digits of 3.14 as 4-bit binary numbers (eg  $3 = 0011, 4 = 0100$  etc), and encode each of them.
  - (c) You receive the following message: 1001011, 0101111, 1101001, 1110010. Correct and decode the message.
  - (d) List all codewords of Hamming's  $[7, 4]$ -code.

2. Let  $C$  be a binary linear code with generator matrix

$$G = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 \end{bmatrix}.$$

Write a parity matrix  $H$  for the code.

3. Write the generator matrix and parity matrix for the ISBN-10 code.
4. Write the generator and parity matrix for the Hamming  $[9, 4]$  square code.
5. (optional) Consider working over the field  $\mathbb{F}_5$ . Which of the following two matrices would be a valid generating matrix  $G$ ? Explain your answer. What problem would one of them cause?

$$G_1 = \begin{bmatrix} 1 & 3 & 0 \\ 3 & 1 & 2 \end{bmatrix}, \quad G_2 = \begin{bmatrix} 1 & 3 \\ 3 & 1 \\ 0 & 2 \end{bmatrix}$$

## Math 512 - Homework 6

Due March 18, 2019

(Additional problems for Homework 7 are added with a \* on them. For probabilities, assume we are transmitting across a noisy memoryless symmetric binary channel with symbol error  $p$ .)

1. Calculate the minimum distance of the Hamming  $[9, 4]$  square code (the one where you input a  $2 \times 2$  matrix and output a  $3 \times 3$  matrix). Since there are only  $2^4 = 16$  codewords, you can just explicitly list them and determine their weights. How many errors can you detect/correct? Write the weight enumerator polynomial.
  - \* What is the probability that a codeword is received with no errors? What is the probability that a codeword, after the error-correction procedure, is the original codeword sent? What is the probability there is an undetected error?
2. Compute the weight enumerator for the  $[n, 1]$ -repetition code. How many errors can you detect? How many errors can you correct?
  - \* What is the probability that a codeword is received with no errors? What is the probability that a codeword, after the error-correction procedure, is the original codeword sent? What is the probability there is an undetected error?
3. The weight enumerator for the  $[4, 2]$ -repetition code is related to the weight enumerator for the  $[2, 1]$ -repetition code. How? Can you guess how this generalizes?
4. Consider the  $[5, 2]$ -code given in class (the one where I handed out the standard array; you can use that table).
  - (a) Correct the message 01010, 11011, 11101.
  - (b) Give an example (or multiple) of introducing 2 errors to a valid codeword, using the standard array to correct, and producing a different codeword than what you started with.
5. Consider the binary  $[6, 3]$  linear code with

$$G = \begin{bmatrix} 1 & 0 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 \end{bmatrix}.$$

We have already determined its minimal distance is  $d = 3$ .

- (a) Construct the parity matrix  $H$ . (You did this in HW 3; you can just copy if you want.)
- (b) Set up a table for *syndrome decoding*. (Note that to do syndrome decoding, you only have to list the coset leaders (elements in row with minimal weight) and their syndrome.)
- (c) Using your table, correct the message 011010, 001110, 100001, 011110, 101111.
  - \* Calculate the weight enumerator polynomial for this code. What is the probability of there being an undetected error? Suppose that, any time we receive a word that doesn't have a unique closest word (i.e. it is not in the "correctable" portion of our syndrome decoding table) we ask for it to be retransmitted. What is the probability a word will have to be retransmitted?

Math 512 - Homework 7  
Due April 1, 2019

1. Show the polynomial  $x^2 + x + 1 \in \mathbb{F}_2[x]$  has no roots by plugging in  $x = 0$  and  $x = 1$ .
2. Perform the following multiplication:  $(x^3 + x^2 + 1)(x + 1) \in \mathbb{F}_2[x]$ .
3. In  $\mathbb{F}_2[x]$ ,  $x^5 + x^2 + x + 1 = (x^2 + 1)g(x)$ . Find  $g(x)$ .
4. In  $\mathbb{F}_2[\alpha]$ , use polynomial division to write  $\alpha^7$  as  $f(\alpha)(\alpha^4 + \alpha + 1) + g(\alpha)$ , where  $\deg(g) \leq 3$ .  
Hint: Compute  $\frac{\alpha^7}{\alpha^4 + \alpha + 1}$

Math 512 - Homework 8  
Due April 8, 2019

1. Let  $\sim$  be the equivalence relation on  $\mathbb{Z}$  defined by

$$x \sim y \quad \text{if } y - x = 2n \text{ for some } n \in \mathbb{Z}.$$

Prove that  $\sim$  is reflexive ( $x \sim x$  for all  $x \in \mathbb{Z}$ ) and symmetric (if  $x \sim y$  then  $y \sim x$ ).

2. Let  $g = g(x) \in \mathbb{Z}[x]$  be a polynomial. Define an equivalence relation  $\sim$  on  $\mathbb{Z}[x]$  by

$$f_1 \sim f_2 \quad \text{if } f_2 - f_1 = g \cdot h \text{ for some } h = h(x) \in \mathbb{Z}[x].$$

Prove that  $\sim$  is transitive. (The proof follows the exact same logic and format as the proof of transitivity we did in class.)

3. Make a chart (as described in class) that calculates  $\alpha^k$  in

$$\mathbb{F}_8 = \mathbb{F}_2[\alpha] / (1 + \alpha + \alpha^3).$$

You can stop when you reach  $\alpha^n = 1$  for some  $n > 1$ .

4. Make a chart (as described in class) that calculates  $\alpha^k$  in

$$\mathbb{F}_{16} = \mathbb{F}_2[\alpha] / (1 + \alpha + \alpha^4).$$

You can stop when you reach  $\alpha^n = 1$  for some  $n > 1$ .

5. Using the model of  $\mathbb{F}_{16}$  from the previous problem, calculate

$$\alpha^4 + \alpha^8 \text{ and } (\alpha^2 + \alpha^3)(1 + \alpha^2).$$



Math 512 - Homework 9  
Due April 15, 2019

1. Consider the generating polynomial  $g(x) = 1 + x^2 + x^3 \in \mathbb{F}_2[x]/(x^7 + 1)$ . This is the example we did in the first half of class (though I didn't write the quotient part until later).
  - (a) Encode the "message"  $1 + x^2$ .
  - (b) By considering the polynomial code as a linear  $[7, 4]$ -code over  $\mathbb{F}_2$  (as we did in class), encode the message 1001.
  - (c) You receive the polynomial message  $x + x^3 + x^4 + x^5$ . Use polynomial division to determine if there is some  $f(x)$  such that  $f(x) \cdot g(x) = x + x^3 + x^4 + x^5$ . Can you "decode" the message?
  - (d) Use polynomial division to find  $h(x)$  such that  $g(x)h(x) = x^7 + 1 \in \mathbb{F}_2[x]$ .

2. Consider the polynomial code over  $\mathbb{F}_8$ , given by the generating polynomial

$$g(x) = (x + \alpha^5)(x + \alpha^6) \in \mathbb{F}_8[x].$$

Here we use the explicit model  $\mathbb{F}_8 := \mathbb{F}_2[\alpha]/(\alpha^3 + \alpha + 1)$ . This determines a  $[7, 5]$ -code over  $\mathbb{F}_8$ , or a  $[21, 15]$ -code over  $\mathbb{F}_2$ . Note: The polynomial  $g(x)$  is different than the one used in class, but the concepts and the code behavior is the same.

- (a) Use  $g(x)$  to encode the following string of 15 bits in a string of 21 bits:

000 010 000 101 100.

- (b) Determine the generator matrix  $G$  of the corresponding linear  $[7, 5]$ -code over  $\mathbb{F}_8$ .
- (c) Determine the generator matrix of the corresponding linear  $[21, 15]$ -code over  $\mathbb{F}_2$ .

Math 512 - Homework 10 (really 11)  
Due April 22, 2019

1. Let  $f(x) \in \mathbb{F}[x]$  and  $a \in \mathbb{F}$ , where  $\mathbb{F}$  is some field (e.g.  $\mathbb{R}, \mathbb{F}_2, \mathbb{Q}$ ; This assures us polynomial division will work well). Prove that  $(x - a)$  divides  $f$  in  $\mathbb{F}[x]$  if and only if  $f(a) = 0$ .  
*Hint: We proved the if direction in class. To show the only if direction, use the fact that the polynomial division algorithm will provide polynomials  $Q(x), R(x) \in \mathbb{F}[x]$  with  $\deg(R) < \deg(x - a)$  such that  $f(x) = Q(x) \cdot (x - a) + R(x)$ .*
2. Consider the polynomial code in  $R_7 = \mathbb{F}_2[x]/(x^7 + 1)$  generated by  $g(x) = x + 1$ .
  - (a) Write a generator matrix  $G$  for the induced cyclic  $[7, 6]$  code.
  - (b) Perform row reduction on the generator matrix  $G$ . (Note that while this changes how one might perform the encoding, the set of all codewords does not change under these row operations. Your row reduced generator matrix is still a generator matrix for the same code.)
  - (c) Show the induced cyclic code is the same as the  $[7, 6]$  code given by adding a parity bit to every codeword.
3. Consider the cyclic code generated by  $g(x) = 1 + x + x^3 \in R_7 = \mathbb{F}_2[x]/(x^7 + 1)$ .
  - (a) Write a generator matrix  $G$  for the induced cyclic  $[7, 4]$  code.
  - (b) Swap the columns in the generator matrix  $G$  in the following way:

$$1 \mapsto 3, \quad 2 \mapsto 1, \quad 3 \mapsto 4, \quad 4 \mapsto 2.$$

In other words, the third column in the new matrix will be the first column in the original matrix. The new generating matrix from the previous part gives a code that is *equivalent, but not equal* to the original code.

- (c) Show this new code is the Hamming  $[7, 4]$  code. You can do this by performing row reduction on the generator matrix from (b) and the generator matrix for the Hamming  $[7, 4]$  code. Conclude that the Hamming code is equivalent to a cyclic code.
4. Consider  $g(x) = x^2 \in \mathbb{F}_2[x]/(x^3 - 1) = R_3$ .
  - (a) Suppose we created a  $[3, 1]$  code  $C$  by using the corresponding generator matrix
$$G = \begin{bmatrix} 0 & 0 & 1 \end{bmatrix}.$$
Show the induced code  $C$  is not cyclic.
  - (b) Show that  $x^2$  does not divide  $x^3 - 1$  in  $\mathbb{F}_2[x]$ .
  - (c) Optional: In fact, show that  $x^2$  is a unit in  $R_3$ , i.e. that  $x^2 \cdot f(x) = 1 \in R_3$  for some  $f(x)$ . Conclude that the ideal generated by  $x^2$  in  $R_3$  is equal to all of  $R_3$ .
  - (d) Does this example contradict the following theorem that was stated in class?  
Assume  $q$  is a power of  $p$ , and  $p \nmid n$ . A linear code  $C \subset \mathbb{F}_q^n$  is cyclic if and only if  $C$  is induced by a generating polynomial  $g(x) \in \mathbb{F}_q[x]$  such that  $g(x)h(x) = x^n - 1$  for some  $h(x) \in \mathbb{F}_q[x]$ .